

# On the Hardness of Learning With Physical Rounding and Noise from Learning With Error

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# Physical security

## Analysis of computational resources

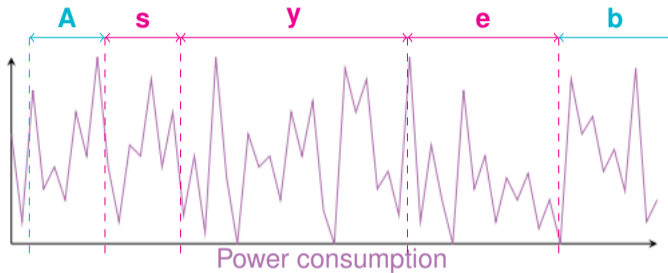
- 1:  $\mathbf{A} \leftarrow \mathbb{Z}_q^{m \times n}$
- 2:  $\mathbf{s} \leftarrow \mathbb{Z}_q^n$
- 3:  $\mathbf{y} \leftarrow \mathbf{A} \cdot \mathbf{s}$
- 4:  $\mathbf{e} \leftarrow \mathcal{D}_\sigma(\mathbb{Z}_q^m)$
- 5:  $\mathbf{b} \leftarrow \mathbf{y} \oplus \mathbf{e}$



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# Physical security

Deduction of information on sensitive data

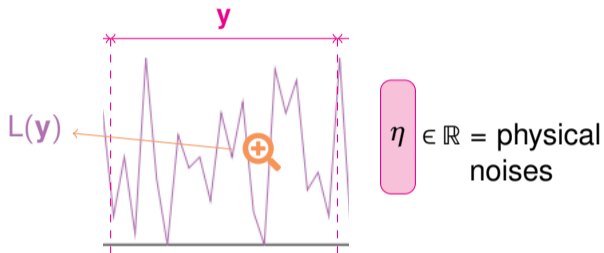
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# Physical security

How much sensible ?

- 1:  $\mathbf{A} \xleftarrow{\$} \mathbb{Z}_q^{m \times n}$
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## Leveraging problems

- $(\mathbf{A}, \mathbf{y}) \rightarrow \text{Not Hard}$
- $(\mathbf{A}, \text{HW}(\mathbf{y})) \rightarrow \text{Hard ?}$
- $(\mathbf{A}, \text{HW}(\mathbf{y}) + \eta) \rightarrow \text{Hard}$

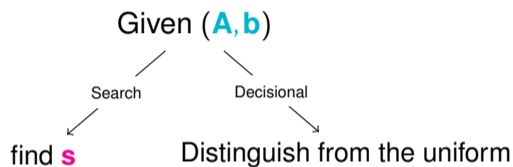
# Learning With Error - LWE

$$\begin{matrix} m & & n & & & & \\ \text{b} & = & \text{A} & \cdot & \text{s} & + & \text{e} & \pmod{q} \\ \mathbb{Z}_q^m & & \mathbb{Z}_q^{m \times n} & & \mathbb{Z}_q^n & & \mathbb{Z}_q^m & \end{matrix}$$

# Learning With Error - LWE

## Variants

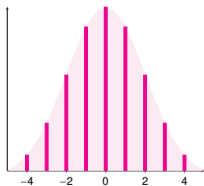
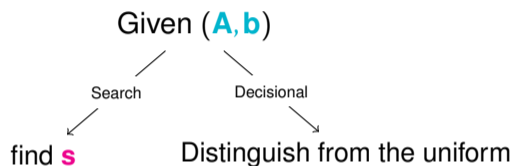
$$\begin{matrix} m \\ \mathbb{Z}_q^m \end{matrix} \mathbf{b} = \begin{matrix} m & n \\ \mathbb{Z}_q^{m \times n} \end{matrix} \mathbf{A} \cdot \begin{matrix} n \\ \mathbb{Z}_q^n \end{matrix} \mathbf{s} + \begin{matrix} m \\ \mathbb{Z}_q^m \end{matrix} \mathbf{e} \pmod{q}$$



# Learning With Error - LWE

## Hardness

$$\underset{\mathbb{Z}_q^m}{\mathbf{b}} = \underset{\mathbb{Z}_q^{m \times n}}{\mathbf{A}} \cdot \underset{\mathbb{Z}_q^n}{\mathbf{s}} + \underset{\mathbb{Z}_q^m}{\mathbf{e}} \pmod{q}$$

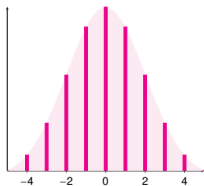
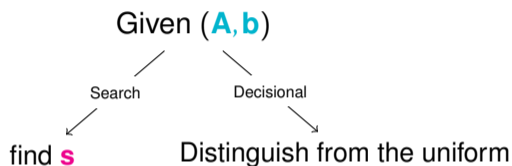


- Security based on worst-case lattices assumption for Gaussian error  $\mathbf{e}$  [Regev05]

# Learning With Error - LWE

## Standardization

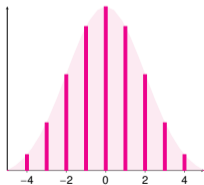
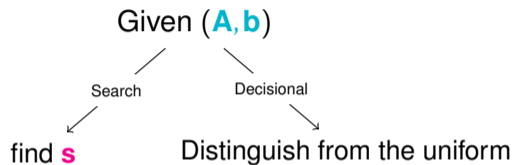
$$\underset{\mathbb{Z}_q^m}{\mathbf{b}} = \underset{\mathbb{Z}_q^{m \times n}}{\mathbf{A}} \cdot \underset{\mathbb{Z}_q^n}{\mathbf{s}} + \underset{\mathbb{Z}_q^m}{\mathbf{e}} \pmod{q}$$



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- Among 5 standards of the NIST, 3 are based on structured LWE variants

# Learning With Error - LWE

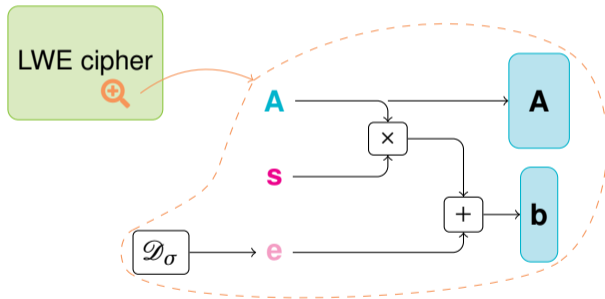
$$\underset{\mathbb{Z}_q^m}{\mathbf{b}} = \underset{\mathbb{Z}_q^{m \times n}}{\mathbf{A}} \cdot \underset{\mathbb{Z}_q^n}{\mathbf{s}} + \underset{\mathbb{Z}_q^m}{\mathbf{e}} \pmod{q}$$



- Security based on worst-case lattices assumption for Gaussian error  $\mathbf{e}$  [Regev05]
- Among 5 standards of the NIST, 3 are based on structured LWE variants

## What about physical attacks ?

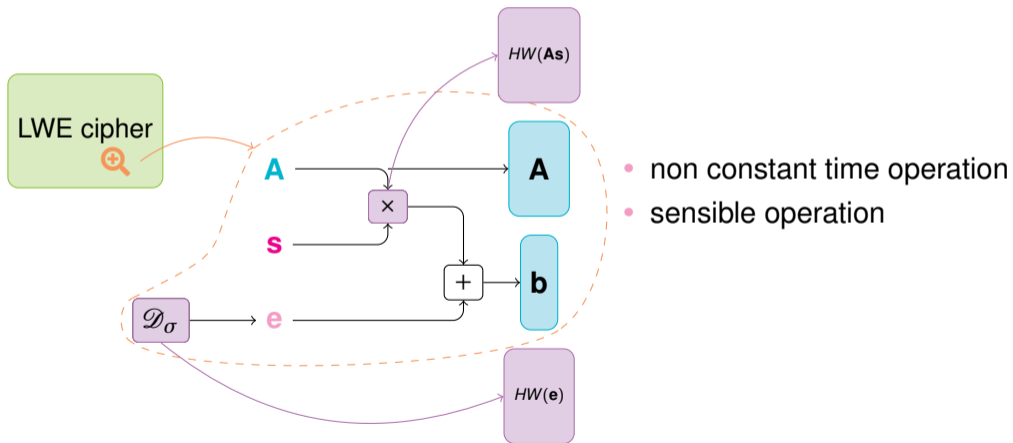
# Side-channel analysis



- power consumption
- timing
- electromagnetic emission
- acoustic emission

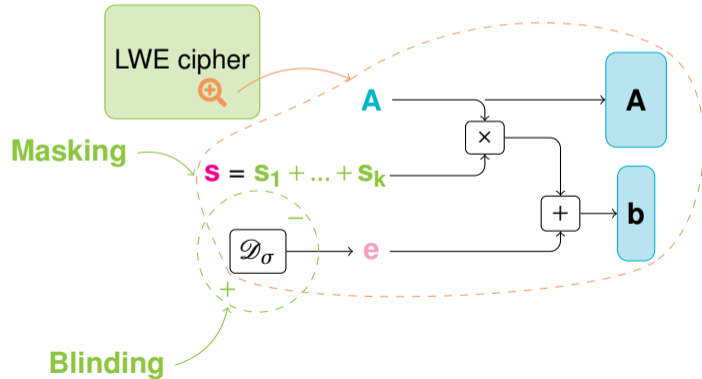
# Side-channel analysis

## Possible leakages



# Side-channel analysis

## Countermeasures

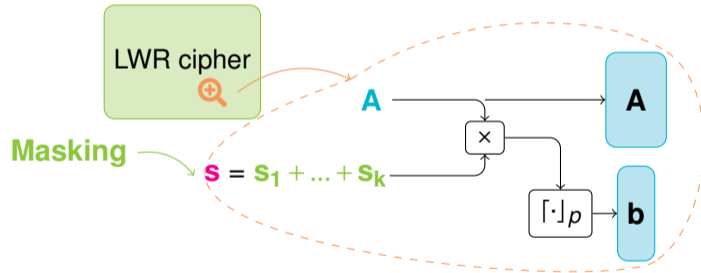


- Masking = share
- Blinding = randomize

**COSTLY**

# Side-channel analysis

Countermeasures: avoiding blinding

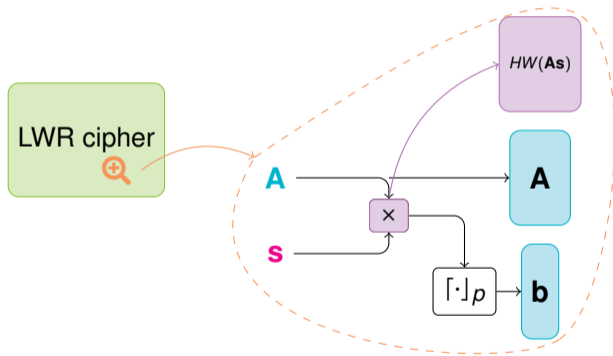


- $\text{LWE: } \mathbf{b} = \mathbf{A}\mathbf{s} + \mathbf{e}$
- $\text{LWR: } \mathbf{b} = \lceil \mathbf{A}\mathbf{s} \rceil_p$  [BPR12]

**New Assumption**

# Side-channel analysis

Countermeasures: avoiding masking



**Relaxed Assumption**

# Hard Learning Problem from Side-channel analysis

LWE Cipher

**A**

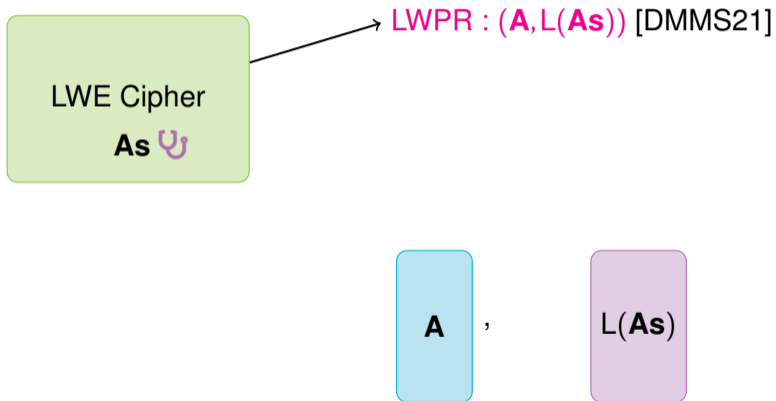
,

**b**

# Hard Learning Problem from Side-channel analysis

Learning With Physical Rounding

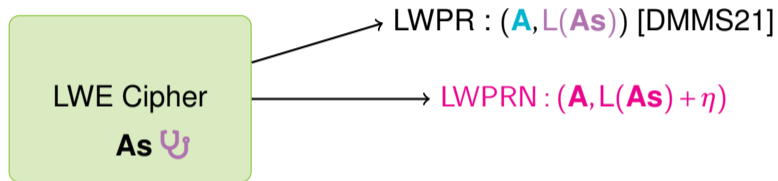
$L = \text{Leakage Function}$



# Hard Learning Problem from Side-channel analysis

Learning With Physical Rounding and Noises

$L = \text{Leakage Function}$



$$\mathbf{e} \in \mathbb{Z}_q^m$$

Mathematical

$$\eta \in \mathbb{R}^m$$

Physical

$$\mathbf{A}$$

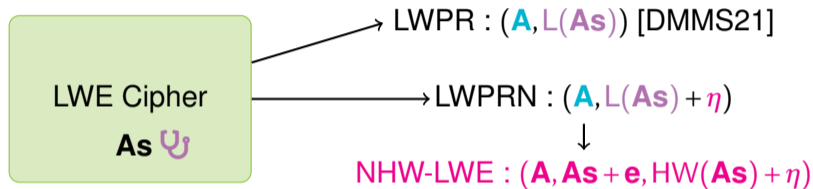
,

$$L(\mathbf{As}) + \eta$$

# Hard Learning Problem from Side-channel analysis

## Noisy Hamming Weight LWE

L = Leakage Function



$$\mathbf{e} \in \mathbb{Z}_q^m$$

Mathematical

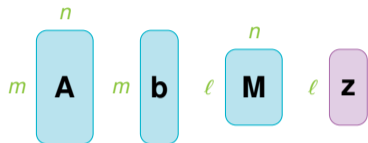
$$\eta \in \mathbb{R}^m$$

Physical

$$\mathbf{A}, \mathbf{b}, \text{HW}(\mathbf{As}) + \eta$$

# LWE variants

## LWE with Hints [KLSS23,LSW25]



$$\begin{matrix} l \\ \mathbf{z} \\ \mathbb{Z}_q^l \end{matrix} = \begin{matrix} l & n \\ \mathbf{M} \\ \mathbb{Z}_q^{l \times n} \end{matrix} \cdot \begin{matrix} n \\ \mathbf{s} \\ \mathbb{Z}_q^n \end{matrix} + \begin{matrix} l \\ \mathbf{f} \\ \mathbb{Z}_q^m \end{matrix} \pmod{q}$$

### Choice of $M$

- Adversary
- (possibly) unstructured

# LWE variants

## LWE with Hints [KLSS23,LSW25]

$$\begin{array}{c} \begin{array}{cccc} \overset{n}{\text{A}} & \text{b} & \overset{n}{\text{M}} & \text{z} \\ \text{m} & \text{m} & \ell & \ell \end{array} \\ \\ \ell \text{ z} = \ell \text{ M} \cdot \text{s} + \text{f} \pmod{q} \\ \mathbb{Z}_q^\ell \quad \mathbb{Z}_q^{\ell \times n} \quad \mathbb{Z}_q^n \quad \mathbb{Z}_q^m \end{array}$$

### Choice of $\mathbf{M}$

- Adversary
- (possibly) unstructured

## Entropic LWE [BD20]

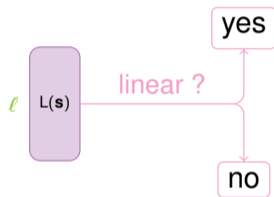
$$\begin{array}{c} \text{b} = \text{A} \cdot \text{s} + \text{e} \pmod{q} \\ \mathbb{Z}_q^m \quad \mathbb{Z}_q^{m \times n} \quad \mathbb{Z}_q^n \quad \mathbb{Z}_q^m \end{array}$$

The diagram shows a dashed circle around the secret  $\mathbf{s}$  and error  $\mathbf{e}$  terms, with a pink circle labeled  $\mathcal{S}$  inside, representing the entropy set.

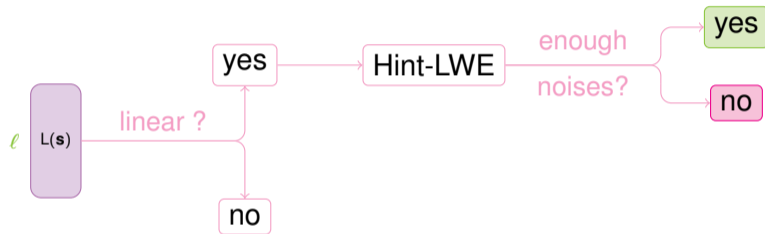
### Choice of $\mathcal{S}$

- Large entropy
- Short secret variant

# Considering additional leakage



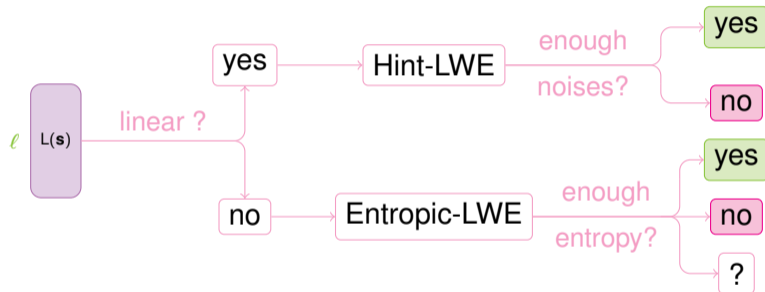
# Considering additional leakage



**Hard Problem**

**Easy Problem**

# Considering additional leakage



**Hard Problem**

**Easy Problem**

# Summary

- We have

LWE Cipher  
As  $\Psi$

+



# Summary

- We have

LWE Cipher  
 $As \cup$

+



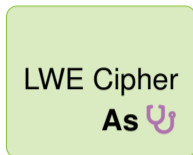
- We want

$A$  ,  $L(As)$  +  $\eta$

LWPRN hard

# Summary

- We have



+



- We want



LWPRN hard

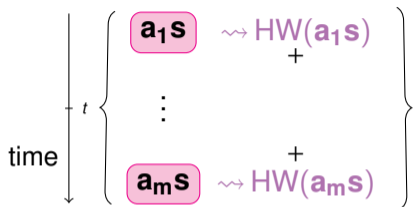
- We study



hardness of NHW-LWE  
under LWE assumption

# Do variants apply to our case ?

Parallel

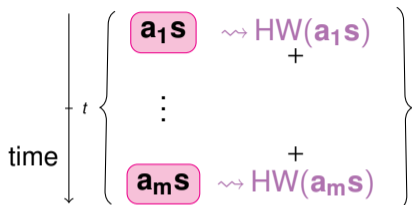


$$\sum_{i=1}^m \text{HW}(\mathbf{a}_i \mathbf{s}) \rightarrow \log(m \log(q))$$

Leakage's  
entropy

# Do variants apply to our case ?

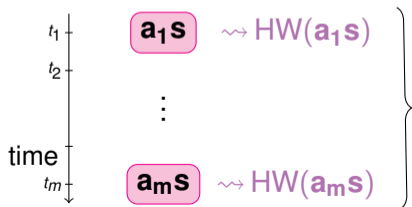
Parallel



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Leakage's  
entropy

Serial

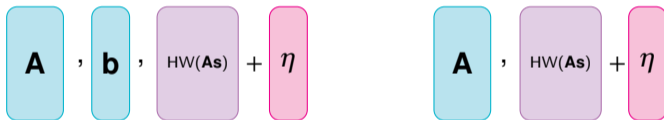


$$(\text{HW}(\mathbf{a}_1 \mathbf{s}), \dots, \text{HW}(\mathbf{a}_m \mathbf{s})) \rightarrow m \log(\log(q))$$

# Results on Noisy Hamming Weight Learning With Error

From NHW-LWE to LWPRN

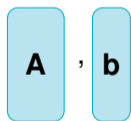
$$\text{NHW-LWE}_{q,n,m,\sigma_e}^{\sigma_\eta} \xrightarrow{\text{Remove } \mathbf{b}} \text{LWPRN}_{q,n,m,\sigma_e}^{\text{HW},\sigma_\eta}$$



# Results on Noisy Hamming Weight Learning With Error

From LWE to NHW-LWE

$$\text{LWE}_{q,n,m,\sigma_e} \xrightarrow{\text{Theorem}} \text{NHW-LWE}_{q,n,m,\sigma_e}^{\sigma_\eta} \xrightarrow{\text{Remove } \mathbf{b}} \text{LWPRN}_{q,n,m,\sigma_e}^{\text{HW},\sigma_\eta}$$



- Statistical distance:  $\frac{\lceil \log_2(\sqrt{\lambda}\sigma_e) \rceil}{\sigma_\eta} \leq \varepsilon$
- Rényi divergence:  $\sigma_\eta \geq \sqrt{m} \lceil \log_2(\sqrt{\lambda}\sigma_e) \rceil$

## Intuition of the reduction

LWE  $\Rightarrow$  NHW-LWE

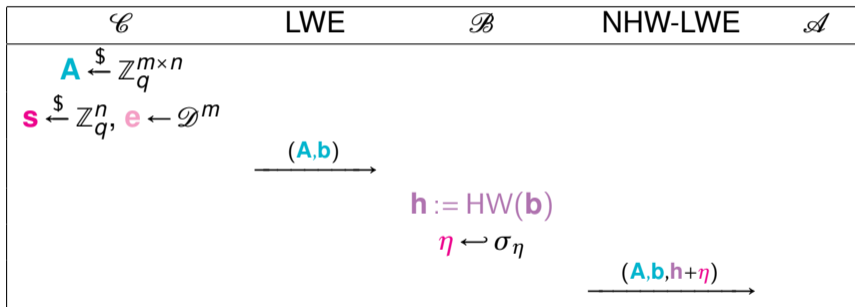
Construct  $\mathcal{B}$  an adversary against LWE using  $\mathcal{A}$  an adversary against NHW-LWE.

# Intuition of the reduction

From LWE sample to NHW-LWE sample

## LWE $\Rightarrow$ NHW-LWE

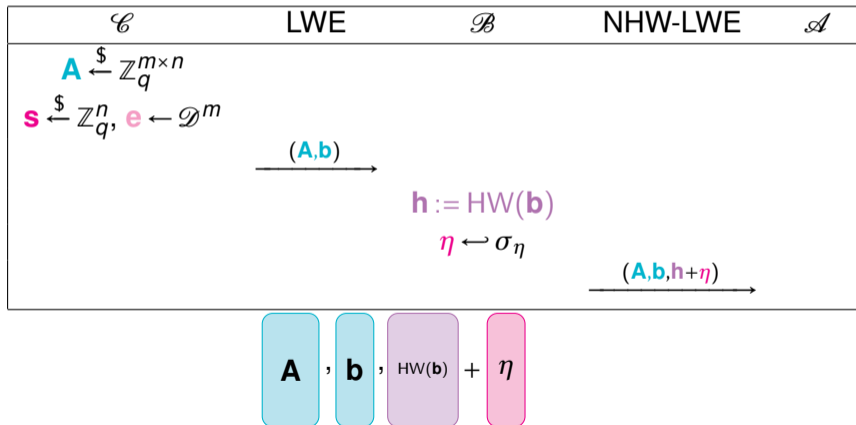
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# Intuition of the reduction

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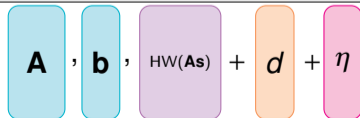
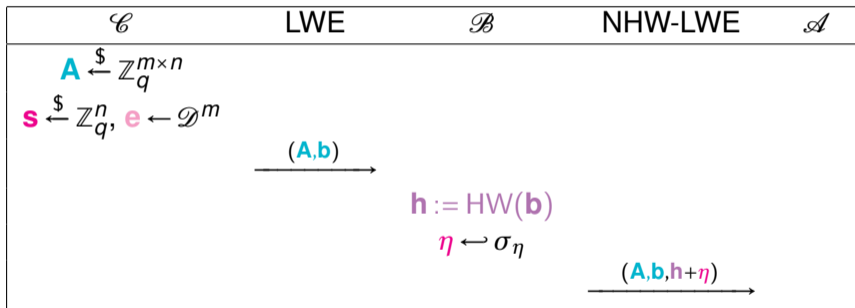


# Intuition of the reduction

Rewriting HW(**b**)

**LWE**  $\Rightarrow$  **NHW-LWE**

Construct  $\mathcal{B}$  an adversary against LWE using  $\mathcal{A}$  an adversary against NHW-LWE.

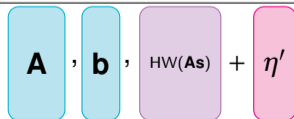
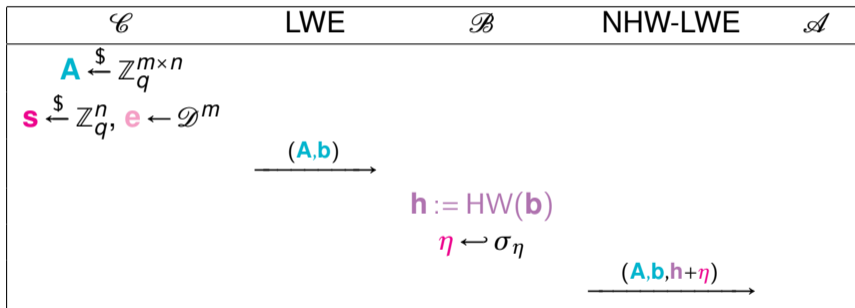


# Intuition of the reduction

Masking  $d$

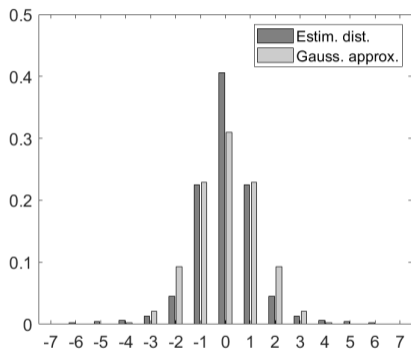
LWE  $\Rightarrow$  NHW-LWE

Construct  $\mathcal{B}$  an adversary against LWE using  $\mathcal{A}$  an adversary against NHW-LWE.

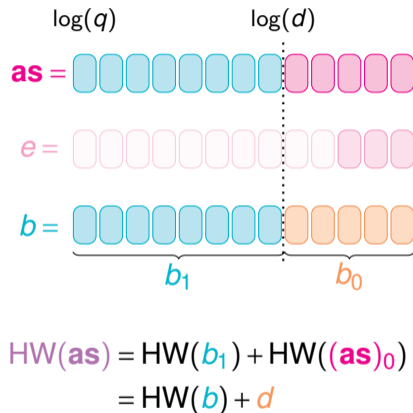


# From LWE sample to NHW-LWE sample

Influence of the LWE error

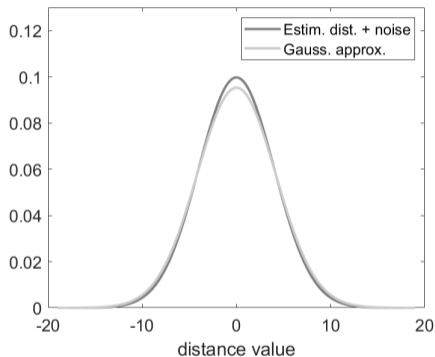


Difference  $\text{HW}(\mathbf{b}) - \text{HW}(\mathbf{As})$  and its Gaussian approximation.



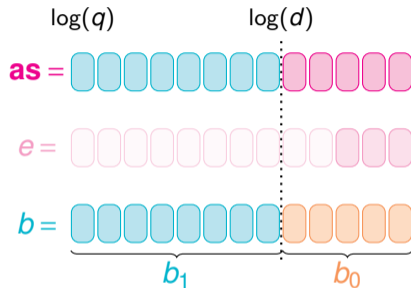
# From LWE sample to NHW-LWE sample

Masking this influence



Difference

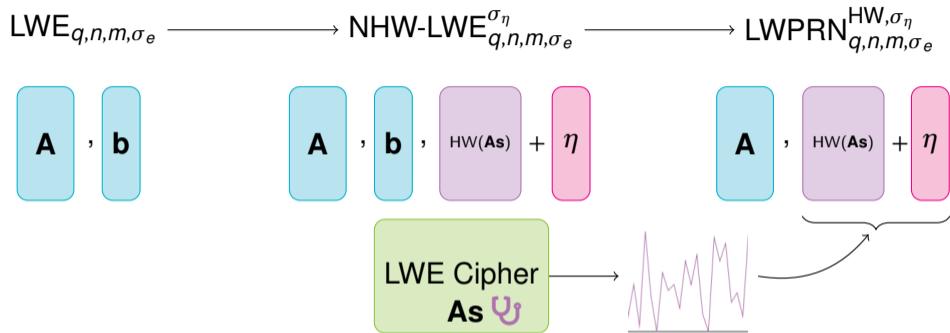
$(HW(\mathbf{b}) + \eta) - (HW(\mathbf{As}) + \eta)$  and its  
Gaussian approximation.



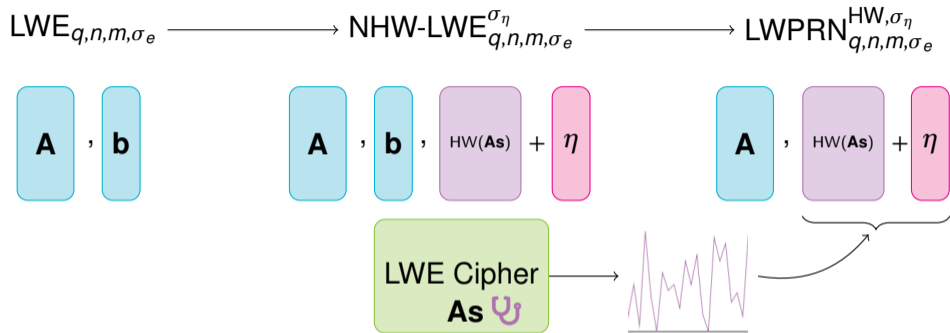
$$\begin{aligned} HW(\mathbf{as}) + \eta &= HW(\mathbf{b}_1) + d + \eta \\ &\approx HW(\mathbf{b}) + \eta \end{aligned}$$

$\eta$  independent of  $q$

# Conclusion and Open Problems

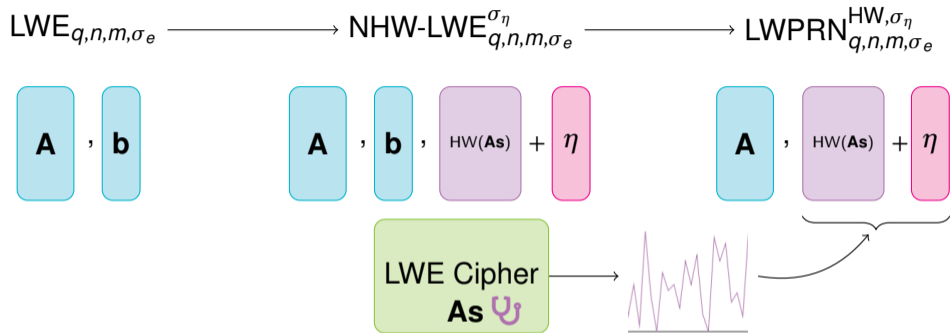


# Conclusion and Open Problems



- How can we better the bound  $d$  knowing  $\mathbf{A}$  ?
- Relaxation on masking ?

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**Thank you for your attention !**